

The edge-count criterion for graphic lists

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Abstract

We give a new short proof of Koren’s characterization of graphic lists, extended to multigraphs with bounded multiplicity p , called p -graphs. The Edge-Count Criterion (ECC) for an integer n -tuple d and integer p is the statement that for all disjoint sets I and J of indices, $\sum_{i \in I} d_i + \sum_{j \in J} [p(n-1) - d_j] \geq p|I||J|$. An integer list d is the degree list of a p -graph if and only if it has even sum and satisfies ECC. Analogous statements hold for bipartite or directed graphs.

The problem of characterizing degree lists (also called “degree sequences”) of simple graphs is well studied. The sum is twice the number of edges and hence must be even, but this condition is not sufficient. Sierksma and Hoogeveen [10] summarized seven characterizations. With additional results, these also appear in [7]. The various characterizations have been proved in many ways; we will not attempt to survey the proofs.

We give a new short proof of another natural characterization, due to Koren [6], which we call the Edge-Count Criterion. Koren used it to characterize the polytope of degree lists [9]. We prove the characterization in the more general setting of multigraphs with bounded multiplicity p . The idea also works for bipartite or directed graphs.

A multigraph G with bounded multiplicity p is a pair consisting of a set $V(G)$ of vertices and a multiset $E(G)$ of unordered pairs of vertices, where each pair occurs at most p times as an edge. Motivated by Berge [1], we call such a multigraph a p -graph (the 1-graphs are the *graphs* or *simple graphs*). Let $\mu(xy)$ denote the multiplicity of an edge xy ; if $\mu(xy) > 0$, then x and y are *adjacent*. The *complement* of a p -graph G , denoted \overline{G} , is the p -graph with vertex set $V(G)$ such that $\mu_{\overline{G}}(xy) = p - \mu_G(xy)$ for all $xy \in \binom{V(G)}{2}$. The *degree* of a vertex v , written $d(v)$, is the sum of the multiplicities of the pairs containing v . We write an integer list (d_1, \dots, d_n) simply as d . An integer n -tuple d is *p -graphic* if the entries are the vertex degrees of some p -graph. Such a p -graph is a *realization* of d . Let $[n] = \{1, \dots, n\}$.

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Definition 1. An integer n -tuple d satisfies the Edge-Count Criterion (ECC) for p -graphs if for all $I, J \subseteq [n]$ with $I \cap J = \emptyset$,

$$\sum_{i \in I} d_i + \sum_{j \in J} [p(n-1) - d_j] \geq p|I||J|. \quad (*)$$

We call this the Edge-Count Criterion because always $\mu_G(xy) + \mu_{\overline{G}}(xy) = p$. The sum on the left counts degrees in G for vertices of I and in \overline{G} for vertices of J . The total must account for the total multiplicity of all pairs in $I \times J$, regardless of how it splits between G and \overline{G} . Thus the condition is necessary. We will give a short proof that when even sum is also required it becomes sufficient.

Koren's statement of the ECC for 1-graphs, when expressed in our notation, was $\sum_{j \in J} d_j \leq \sum_{i \in I} d_i + |J|(n-1-|I|)$. We have reordered the terms to facilitate a short proof and express the natural generalization to p -graphs. Characterizations of p -graphic lists were given by Chungphaisan [2] and by Berge [1].

Fulkerson–Hoffman–McAndrew [5] proved that every 1-graphic list has a realization in which any specified vertex v is adjacent to vertices whose degrees are the largest entries in the list other than its own. We need the extension to p -graphs of an easy special case.

Lemma 2. Let d be a p -graphic list with largest entry k . If $d_j > 0$ and d_j is not the only k in d , then in some realization of d a vertex of degree d_j is adjacent to a vertex of degree k .

Proof. Let G be a realization of d . Let x and v be vertices of degrees k and d_j . If $\mu(xv) = 0$, then v is adjacent to some other vertex u . Since $d(u) \leq k$, and v is adjacent to u but not x , there exists y such that $\mu(xy) > \mu(uy)$. Decreasing $\mu(xy)$ and $\mu(vu)$ by 1 and increasing $\mu(xv)$ and $\mu(uy)$ by 1 yields a realization as desired. \square

Theorem 3. An integer n -tuple d with even sum is p -graphic if and only if it satisfies the ECC for p -graphs.

Proof. We have observed that the conditions are necessary. For sufficiency, we use induction on $n + \sum d_i$. For a list d , let $S_{I,J}(d)$ denote $\sum_{i \in I} d_i + \sum_{j \in J} [p(n-1) - d_j]$, so ECC states that $S_{I,J}(d) \geq p|I||J|$ whenever I and J are disjoint.

Suppose that d satisfies ECC. Using pairs I, J in which one set is empty and the other is a singleton, we obtain $0 \leq d_i \leq p(n-1)$ for all i , so the induction parameter is positive. When it equals 1, the unique realization is the unique 1-vertex p -graph, which has no edges.

For the induction step, index d so that d_1 is a largest entry and d_n is smallest. If $d_n = 0$, then form d' by deleting d_n . Since d' is an $(n-1)$ -tuple and ECC holds for d , we have $S_{I,J}(d') = S_{I \cup \{n\}, J}(d) - p|J| \geq p(|I|+1)|J| - p|J| = p|I||J|$. Thus ECC holds for d' , which has the same sum as d . By the induction hypothesis d' is p -graphic, and adding an isolated vertex to a realization of d' yields a realization of d .

Hence we may assume $d_n \geq 1$. Form d' by subtracting 1 from the first and last entries. If d' is p -graphic, then applying Lemma 2 to the complement of a realization of d' yields a realization of d' having vertices x and y of degrees $d_1 - 1$ and $d_n - 1$ such that $\mu(xy) < p$. Increasing the multiplicity of xy completes a realization of d .

Since d' has even sum, by the induction hypothesis it suffices to show that d' satisfies ECC. If $d'_i > d'_j$ for some $i \in I$ and $j \in J$, then moving i to J and j to I reduces $S_{I,J}(d')$ without changing $|I||J|$. Hence it suffices to prove (*) when $d'_i \leq d'_j$ for all $i \in I$ and $j \in J$.

Writing (*) as $\sum_{i \in I} (d'_i - p|J|) + \sum_{j \in J} [p(n-1) - d'_j] \geq 0$, we need only prove (*) when $d'_i < p|J|$ for $i \in I$. Furthermore, if $d'_{i'} < d'_i < p|J|$ and $i \in I$, then $i' \notin J$, and adding i' to I if not already in I reduces the left side. Hence we may assume that all entries smaller than any indexed by I are also indexed by I . Similarly, to ensure $\sum_{i \in I} d'_i + \sum_{j \in J} [p(n-1-|I|) - d'_j] \geq 0$, we may assume that $d'_j > p(n-1-|I|)$ for $j \in J$, and entries larger than any indexed by J are indexed by J .

Since $0 \leq d'_i \leq p(n-1)$, (*) holds when I or J is empty. Hence we may assume that both are nonempty, with J containing the index of a largest entry and I containing that of a smallest. In particular, $n \in I$. If $d'_j = d_1 - 1$ for any $j \in J$ (including $j = 1$), then $S_{I,J}(d') = S_{I,J-\{j\}+\{1\}}(d) \geq |I||J|$. Hence we may assume that $d'_j = d_1$ for $j \in J$.

For $j \in J$, we have $d_1 = d'_j > p(n-1-|I|)$, or $-p(n-1) - d_1 > -p|I|$. If $1 \notin I$, then

$$S_{I,J}(d') = S_{I,J \cup \{1\}}(d) - 1 - [p(n-1) - d_1] \geq p|I|(|J|+1) - p|I| = |I||J|.$$

Hence we may assume $1 \in I$. Now $d'_1 < p|J|$, so $d_1 \leq p|J|$. With $|I| + |J| \leq n$, we have

$$\begin{aligned} S_{I,J}(d') &= \left(\sum_{i \in I} d_i \right) - 2 + [p(n-1) - d_1]|J| \geq d_1 + |I| - 1 - 2 + [p(|I| + |J| - 1) - d_1]|J| \\ &= |I| - 3 + p|I||J| + (|J| - 1)(p|J| - d_1). \end{aligned}$$

Failure requires $(|J| - 1)(p|J| - d_1) < 3 - |I|$ and equality throughout the computation. Hence $I = \{1, n\}$ and $|J| \in \{1, d_1/p\}$; also $d_n = 1$ and $|I| + |J| = n$, so $|J| = n - 2$. If $|J| = 1$, then $d = (d_1, d_1, 1)$. If $p|J| = d_1$, then $d = (p(n-2), \dots, p(n-2), 1)$, with $n-1$ entries equal to $p(n-2)$. In each case, d has odd sum, so these possibilities are excluded.

Hence d' satisfies ECC, and the induction hypothesis applies to complete the proof. \square

When $p = 1$, some cases disappear earlier. The requirement for $d_i = d_j$ with $i \in I$ and $j \in J$ is $p|J| > d'_i = d'_j > p(n-1-|I|)$, which simplifies to $|I| + |J| \geq n - 1 + 2/p$ and cannot hold when $p = 1$. Therefore, when $p = 1$ we may assume that $I = \{i: d'_i < |J|\}$ and $J = \{j: d'_j > n - 1 - |I|\}$. This leads more quickly to $n \in I$ and $d_1 = n - |I|$.

For bipartite graphs there is a similar characterization. A pair of lists (r_1, \dots, r_m) and (s_1, \dots, s_n) is *bigraphic* if there is a bipartite graph with partite sets X and Y such that r

is the list of degrees of vertices in X and s is the list of degrees of vertices in Y . As above, we consider bipartite p -graphs. A characterization follows from a more general result of Ore [8], which we state in our notation: A bipartite graph G with partite sets $[m]$ and $[n]$ has a spanning subgraph with degree lists r for $[m]$ and s for $[n]$ if and only if, whenever $I \subseteq [m]$ and $J \subseteq [n]$, $\sum_{i \in J} s_j$ is at most $\sum_{i \in I} r_i$ plus the number of edges joining J and $[m] - I$. When G is a complete bipartite p -graph, this reduces to the following statement.

Theorem 4. *Integer lists r and s form the degree lists for a bipartite p -graph if and only if they satisfy the Bipartite Edge-Count Criterion that for all $I \subseteq [m]$ and $J \subseteq [n]$,*

$$\sum_{i \in I} r_i + \sum_{j \in J} (pm - s_j) \geq p|I||J|.$$

For $p = 1$, this is known as the Gale–Ryser Theorem. It can be proved using network flow methods or by a short inductive proof. A proof parallel to that of Theorem 3 is also quite short, since the difficult case ($1 \in I$) does not occur in the bipartite setting.

We omit the analogous statement for directed graphs.

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