Distinguishing number of line graphs of complete bipartite graphs

Or: Distinguishing number > 2

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1 2 1 2
2 1 1 2
2 1 2 1
1 1 1 2
2 2 1 1
2 1 2 2
```

Bottom array is an identity coloring - only automorphism is the trivial one

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Determine the smallest c such that an $s \times t$ grid has an identity coloring



Determine the smallest c such that an $s \times t$ grid has an identity coloring - fewest colors to destroy symmetry

Determine the smallest t such that every c coloring of $s \times t$ grid has a non trivial automorphism - smallest 'size' with unavoidable symmetry

Determine the smallest c such that an $s \times t$ grid has an identity coloring

Determine the distinguishing number of the cartesian product of two complete graphs: $D(K_s \square K_t) = ?$ Fewest colors on vertices such that only trivial automorphism

Determine fewest colors on edges of complete bipartite graph such that only trivial automorphism

Determine the distinguishing number of the action of the group $S_s \times S_t$ on the grid $\mathbb{N}_s \times \mathbb{N}_t$

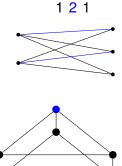
Determine the smallest c such that an $s \times t$ grid has an identity coloring

Determine the distinguishing number of the cartesian product of two complete graphs: $D(K_s \square K_t) = ?$ Fewest colors on vertices such that only trivial automorphism Imrich, Jerebic, Klavzar (2008) formula for most values, recursion for a few

Determine fewest colors on edges of complete bipartite graph such that only trivial automorphism

Fisher and Isaak (2008) formula for most values, recursion for a few

Determine the distinguishing number of the action of the group $S_s \times S_t$ on the grid $\mathbb{N}_s \times \mathbb{N}_t$ Chan (2006) recursion for values



Identity coloring of $s \times t$ grid

- \Leftrightarrow identity edge coloring of $K_{s,t}$ (bipartite adjacency matrix)
- \Leftrightarrow identity vertex coloring of $K_s \square K_t$ (Imrich and Miller result on

automorphisms of $G\Box H$)



Identity 2-coloring of 2 \times 2 grid: $\begin{array}{c} 1 & 1 \\ 1 & 2 \end{array}$

Identity 2-coloring of 3 \times 3 grid: 1 1 1 1 2 1 2 2

Identity 2-coloring of 2
$$\times$$
 2 grid: $\begin{bmatrix} 1 & 1 \\ 1 & 2 \end{bmatrix}$
Identity 2-coloring of 3 \times 3 grid: $\begin{bmatrix} 1 & 1 & 1 \\ 1 & 1 & 2 \\ 1 & 2 & 2 \end{bmatrix}$

Have automorphisms if we allow transposes

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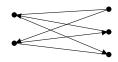
- Have automorphisms if we allow transposes
- Need 3 colors if we allow transposes

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- Have automorphisms if we allow transposes
- Need 3 colors if we allow transposes
- ▶ Otherwise s = t cases we can 'ignore' transposes

Identity orientations

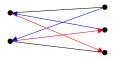
Orient some of the edges of $G \to \text{only trivial automorphism}$ 'Destroying symmetry by orienting edges: complete graphs and complete bigraphs' (Harary and Jacobson 2001) also: Harary and Robinson (2001) and Harary and Ranjan (2005) For bipartite graphs



Identity orientations

Orient some of the edges of $G \rightarrow$ only trivial automorphism 'Destroying symmetry by orienting edges: complete graphs and complete bigraphs' (Harary and Jacobson 2001) also: Harary and Robinson (2001) and Harary and Ranjan (2005)

For bipartite graphs = 3 colors



Cartesian products and distinguishing number

 $G \square H$ - replace vertices of G with copies of H, edges of G become edges between copies of the same vertex of H

G is prime if $G = G_1 \square G_2$ only trivially

$$G\square G\square G=G^3$$
 etc

- ► Hypercubes $D(K_2^n) = 2$ except $D(K_2^2) = D(K_2^3) = 3$ (Bogstead and Cowen 2004)
- Connected prime graphs D(G^r) = 2 for r ≥ 4 (Albertson 2005)
- ► Connected graphs $D(G^r) = 2$ for $r \ge 3$ (Klavzar and Zhu 2007)
- ▶ $D(G^r) = 2$ for $r \ge 2$ except $D(K_2^2) = D(K_2^3) = D(K_3^2) = 3$ (Imrich and Klavzar 2006)

G, H connected relatively prime graphs with $|G| \neq |H|$

$$\Rightarrow D(G\Box H) \leq D(K_{|G|}\Box K_{|H|})$$

automorphisms of $G\square H$ are also automorphisms of $K_{|G|}\square K_{|H|}$

(Imrich, Jerebic and Klavzar 2008)

Bounds on distinguishing number of Cartesian products of complete graphs are general bounds

Columns must be distinct else permute duplicates

```
1 2 1 1 2
1 1 1 2 1 permute columns 2 and 5
1 2 2 2 2
```

 $\Rightarrow t \leq 2^3$ (else duplicated column)

Columns must be distinct else permute duplicates
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$$\Rightarrow t \leq 2^3$$
 (else duplicated column)

Any row permutation $\ \Rightarrow$ grid with all possible columns

$$\Rightarrow t \leq 2^3 - 1$$

►
$$t = 2^3 - 1$$
?
1 1 1 1 2 2 2 2
1 1 2 1 1 2 2
1 2 1 2 1 2 1 2

- Omit any column apply row permutation that preserves the column
 - ⇒ grid with all columns except omitted column
 - ⇒ some column permutation gives original grid

$$\Rightarrow t \leq 2^3 - 2$$

 $t = 2^3 - 2?$ 1 1 2 2 2 2
1 2 1 1 2 2
1 1 1 2 1 2

$$t = 2^3 - 2?$$

$$112222$$

$$121122$$

$$111212$$

row 'degrees' distinct ⇒ only trivial permutations columns distinct

$$\Rightarrow t = 2^3 - 2$$
 has an identity 2-coloring

Switching lemma example

5 colors and 7 rows: Complete to all 5⁷ possible columns

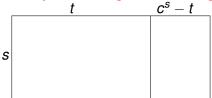
$$\frac{t}{7}$$

- Non-trivial automorphism on right ⇔ Non-trivial automorphism on right
- Switching Lemma: s x t grid has an identity c coloring if and only if s x c^s − t grid does.



- ▶ If *t* is 'small' use 'distinct row degrees', distinct columns construction
- ▶ If t is 'big' (close to c^s) use induction and switching lemma and look for identity coloring of $s \times c^s t$ grid





- ▶ YES if $t \le c^s \lceil \log_c s \rceil 1$
- ▶ NO if $t \ge c^s \lceil \log_c s \rceil + 1$
- ▶ SOMETIMES if $t = c^s \lceil \log_c s \rceil$

$$c^{s}-t$$

- ▶ YES if $t \le c^s \lceil \log_c s \rceil 1$
- ▶ NO if $t \ge c^s \lceil \log_c s \rceil + 1$
- ▶ SOMETIMES if $t = c^s \lceil \log_c s \rceil$
 - ▶ NO if $s \ge c^{1+x} \lfloor \log_c x \rfloor$
 - ▶ YES if $s \le c^{1+x} \lfloor \log_c x \rfloor 2$
 - ▶ SOMETIMES if $s = c^{1+x} \lfloor \log_c x \rfloor 2$
 - $x = \lfloor \log_c(s-1) \rfloor$

Translation to distinguishing number for $s \le t$:

Let
$$c = \lceil (t+1)^{1/s} \rceil$$

smallest c such that $t < c^s$
 $D(K_s \square K_t) =$

$$\begin{cases} c+1 \text{ if } c^s - \lceil \log_c s \rceil + 1 \le t \le c^s - 1 \\ c \text{ if } s \le t \le c^s - \lceil \log_c s \rceil \\ c \text{ or } c+1 \text{ when } t = c^s - \lceil \log_c s \rceil \text{ (determine by recursion)} \end{cases}$$

It will be c if t is 'big' relative to s; e.g. $t \ge s^s$

Grids in higher dimensions? Example: Note that $13 \cdot 42 = 546$

▶ $13 \times 42 \times (c^{546} - \lceil \log_c 546 \rceil - 1)$ grid has an identity c coloring

```
(identity c coloring of 546 \times (c^{546} - \lceil \log_c 546 \rceil - 1) grid)
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General case replace 13,42 with s, t



Summary - we can 'most of the time' directly determine the distinguishing number of $K_s \square k_t$ and otherwise can determine with a recursion using at most iterated logarithm (base c of t) steps. Extends to higher dimensional grids