# Min-max Theorems for the $k$-Path Partition Problem 

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Cumberland Conference May 2014

## Hamiltonian Path



## Hamiltonian Path



## 1-HP (Hamiltonian Path) problem:

 Hamiltonian path with a specified end?

Hamiltonian path


No 1-Hamiltonian path
k-PP (Path Partition/Path Cover) problem: Minimum path partition with $k$ specified ends


3 paths to partition with specified ends


1 path to partition with specified ends


2 paths to partition with specified ends


3 paths to partition with specified ends

## (Regular) Path Partition

- Efficient algorithms and min-max theorems for Co-comparability graphs (threshold, cographs, interval, ...)
- Efficient algorithms but no nice min-max theorem for block graphs (Hamiltonian path even for distance hereditary)


## (with specified ends) k-Path Partition

- Efficient algorithms for
- Our goal
min-max theorems/certificate/certifying algorithm
for 'nice' classes


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## (with specified ends) $k$-Path Partition

- Efficient algorithms for
- Unit interval (Asdre, Nikopoulos 2009; Mertzios, Unger 2010)
- Cographs (includes threshold) (Hung; Asdre, Nikopoulos 2006)
- Block graphs (from regular partition)
- Interval (1-HP only) (Asdre, Nikopoulos 2009)
- Distance hereditary (2-HP only, from regular partition)
min-max theorems/certificate/certifying algorithm


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## $k$-Path Partition ${ }^{\prime}={ }^{\prime}$ (regular) PP with pendants



2-HP
Regular HP
So just add pendants and translate known results:

- works for trees
- Fails for block graphs - no certificate for regular PP
- Fails for other classes - Adding pendant $\Rightarrow$ out of class


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## Certificate for (Regular) Path Partition



## Theorem

If $G$ is (interval graph, threshold graph, ....) then Min Path Partition $=\operatorname{Max} C(G-U)-|U|$

Would like something similar for $k$-Path Partition

## Certificate (gen 1) for Path Partition with terminals $T$



## Theorem

If $G$ is a tree with terminals $T$ then
Min $|T|$-Path Partition $=\operatorname{Max} C(G-U)-|U-T|$
prove directly via induction or
use pendant construction and certificate for regular partitions

## Trivial Certificate for Path Partition with terminals $T$



$$
\text { need } \geq\left\lceil\frac{5}{2}\right\rceil=3 \text { paths }
$$

## Theorem

If $G$ is 2 -connected unit interval graph with terminals $T$ then $\operatorname{Min}|T|$-Path Partition $=\left\lceil\frac{|T|}{2}\right\rceil$ or +1 when ....
$k=2$ case Mertzios and Unger 2010
Note: 2-connected unit interval $\Rightarrow$ Hamiltonian cycle; connected unit interval $\Rightarrow$ Hamiltonian path

## Theorem

If $G$ is 2-connected unit interval graph with terminals $T$ then $\operatorname{Min}|T|$-Path Partition $=\left\lceil\frac{|T|}{2}\right\rceil$ or +1 when (see below)

Unit interval $=$ intersection graph of unit intervals


## Certificate (gen 2) for Path Partition with terminals $T$



## Theorem

If $G$ is a threshold graph with terminals $T$ then
$\operatorname{Min}|T|$-Path Partition $=\operatorname{Max}\left(\sum\left\lceil\frac{\left|T_{i}\right|}{2}\right\rceil\right)+|R|-|U-T|$ or +1 when ....

## Theorem

If $G$ is a threshold graph with terminals $T$ then
$\operatorname{Min}|T|$-Path Partition $=\operatorname{Max}\left(\sum\left\lceil\frac{\left|T_{i}\right|}{2}\right\rceil\right)+|R|-|U-T|$ or +1 when ....


Trees with terminals $T$ :
$\operatorname{Min}|T|$-Path Partition $=\operatorname{Max} C(G-U)-|U-T|$

2-connected unit interval graphs with terminals $T$ :
$\operatorname{Min}|T|$-Path Partition $=\left\lceil\frac{|T|}{2}\right\rceil$
or +1 when ....

Threshold graphs with terminals $T$ :
$\operatorname{Min}|T|$-Path Partition $=\operatorname{Max}\left(\sum\left\lceil\frac{\left|T_{i}\right|}{2}\right\rceil\right)+|R|-|U-T|$ or +1 when ....

Block graphs, unit interval graphs with cut vertices, terminals $T$ : complicated $\min =\max \ldots$.

## Linear Block Graphs (and unit interval with cut vertices)



Impose left right ordering and use 'forced ends'
Decompose along carefully selected cut vertices
i.e., those vertices do not pass through in left/right algorithm


Count the number of 'right' ends in each part
General block graphs: Similar idea with 'tops of paths' and parity of components 'below' special cut vertices

## Block Graphs

Decompose on certain cut vertices count using 'forced ends' and ...


Trees with terminals $T$ :
$\operatorname{Min}|T|$-Path Partition $=\operatorname{Max} C(G-U)-|U-T|$

2-connected unit interval graphs with terminals $T$ :
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