# Min-max Theorems for the *k*-Path Partition Problem

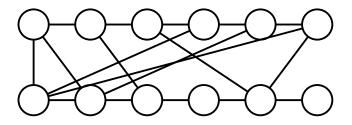
Breeanne Baker The Citidel Garth Isaak Lehigh University

Cumberland Conference May 2014

◆□▶ ◆□▶ ◆□▶ ◆□▶ □□ - のへぐ

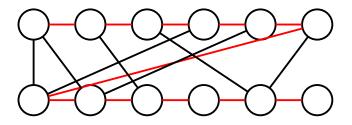
# Hamiltonian Path

◆□▶ ◆□▶ ◆臣▶ ◆臣▶ 三臣 - のへぐ



# Hamiltonian Path

◆□▶ ◆□▶ ◆臣▶ ◆臣▶ 三臣 - のへぐ



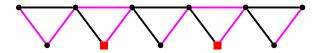
### 1-HP (Hamiltonian Path) problem: Hamiltonian path with a specified end?



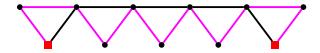
Hamiltonian path

No 1-Hamiltonian path

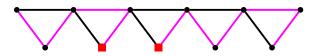
*k*-PP (Path Partition/Path Cover) problem: Minimum path partition with *k* specified ends



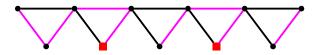
3 paths to partition with specified ends



1 path to partition with specified ends



2 paths to partition with specified ends



3 paths to partition with specified ends

#### (Regular) Path Partition

- Efficient algorithms and min-max theorems for Co-comparability graphs (threshold, cographs, interval, ...)
- Efficient algorithms but no nice min-max theorem for block graphs (Hamiltonian path even for distance hereditary)

#### (with specified ends) k-Path Partition

- Efficient algorithms for
  - Unit interval (Asdre, Nikopoulos 2009; Mertzios, Unger 2010)
  - Cographs (includes threshold) (Hung; Asdre, Nikopoulos 2006)
  - Block graphs (from regular partition)
  - Interval (1-HP only) (Asdre, Nikopoulos 2009)
  - Distance hereditary (2-HP only, from regular partition)
- Our goal -

min-max theorems/certificate/certifying algorithm for 'nice' classes

#### (Regular) Path Partition

- Efficient algorithms and min-max theorems for Co-comparability graphs (threshold, cographs, interval, ...)
- Efficient algorithms but no nice min-max theorem for block graphs (Hamiltonian path even for distance hereditary)

(with specified ends) k-Path Partition

- Efficient algorithms for
  - Unit interval (Asdre, Nikopoulos 2009; Mertzios, Unger 2010)
  - Cographs (includes threshold) (Hung; Asdre, Nikopoulos 2006)
  - Block graphs (from regular partition)
  - Interval (1-HP only) (Asdre, Nikopoulos 2009)
  - Distance hereditary (2-HP only, from regular partition)
- Our goal -

min-max theorems/certificate/certifying algorithm for 'nice' classes

#### (Regular) Path Partition

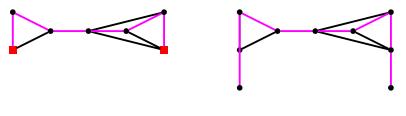
- Efficient algorithms and min-max theorems for Co-comparability graphs (threshold, cographs, interval, ...)
- Efficient algorithms but no nice min-max theorem for block graphs (Hamiltonian path even for distance hereditary)

(with specified ends) k-Path Partition

- Efficient algorithms for
  - Unit interval (Asdre, Nikopoulos 2009; Mertzios, Unger 2010)
  - Cographs (includes threshold) (Hung; Asdre, Nikopoulos 2006)
  - Block graphs (from regular partition)
  - Interval (1-HP only) (Asdre, Nikopoulos 2009)
  - Distance hereditary (2-HP only, from regular partition)
- Our goal -

min-max theorems/certificate/certifying algorithm for 'nice' classes

# k-Path Partition ' = ' (regular) PP with pendants



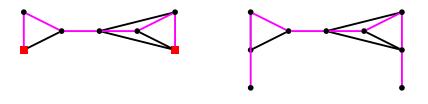
#### 2-HP

Regular HP

So just add pendants and translate known results:

- works for trees
- Fails for block graphs no certificate for regular PP
- Fails for other classes Adding pendant  $\Rightarrow$  out of class

# k-Path Partition ' = ' (regular) PP with pendants



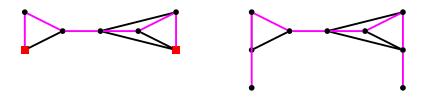
#### 2-HP

Regular HP

# So just add pendants and translate known results:

- works for trees
- Fails for block graphs no certificate for regular PP
- Fails for other classes Adding pendant  $\Rightarrow$  out of class

# k-Path Partition ' = ' (regular) PP with pendants



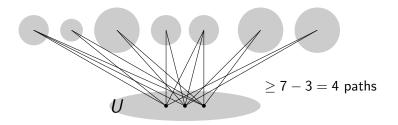
2-HP

Regular HP

So just add pendants and translate known results:

- works for trees
- Fails for block graphs no certificate for regular PP
- Fails for other classes Adding pendant  $\Rightarrow$  out of class

#### Certificate for (Regular) Path Partition

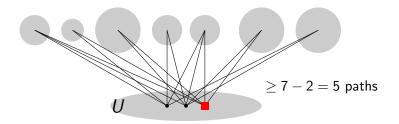


#### Theorem

If G is (interval graph, threshold graph, ....) then Min Path Partition = Max C(G - U) - |U|

#### Would like something similar for k-Path Partition

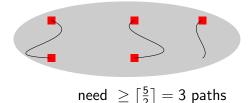
Certificate (gen 1) for Path Partition with terminals T



### Theorem If G is a tree with terminals T then Min |T|-Path Partition = Max C(G - U) - |U - T|

prove directly via induction or use pendant construction and certificate for regular partitions

Trivial Certificate for Path Partition with terminals T



#### Theorem

If G is 2-connected unit interval graph with terminals T then Min |T|-Path Partition  $= \left\lceil \frac{|T|}{2} \right\rceil$ or +1 when ....

・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・

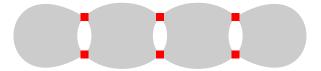
k = 2 case Mertzios and Unger 2010 Note: 2-connected unit interval  $\Rightarrow$  Hamiltonian cycle; connected unit interval  $\Rightarrow$  Hamiltonian path

#### Theorem

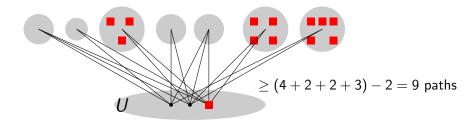
If G is 2-connected unit interval graph with terminals T then Min |T|-Path Partition  $= \left\lceil \frac{|T|}{2} \right\rceil$ or +1 when (see below)

< □ > < 同 > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > <

Unit interval = intersection graph of unit intervals



#### Certificate (gen 2) for Path Partition with terminals T



#### Theorem

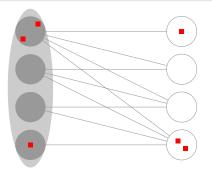
If G is a threshold graph with terminals T then

Min 
$$|T|$$
-Path Partition = Max  $\left(\sum_{i=1}^{|T_i|} | \right) + |R| - |U - T|$   
or +1 when ....

▲□▶ ▲□▶ ▲□▶ ▲□▶ ▲□ ● ● ●

#### Theorem

# If G is a threshold graph with terminals T then Min |T|-Path Partition = Max $\left(\sum \left\lceil \frac{|T_i|}{2} \right\rceil \right) + |R| - |U - T|$ or +1 when ....



・ロト ・雪ト ・ヨト ・ヨー うへぐ

Trees with terminals T: Min |T|-Path Partition = Max C(G - U) - |U - T|

2-connected unit interval graphs with terminals T: Min |T|-Path Partition =  $\left\lceil \frac{|T|}{2} \right\rceil$ or +1 when ....

Threshold graphs with terminals T: Min |T|-Path Partition = Max  $\left(\sum \left\lceil \frac{|T_i|}{2} \right\rceil \right) + |R| - |U - T|$ or +1 when ....

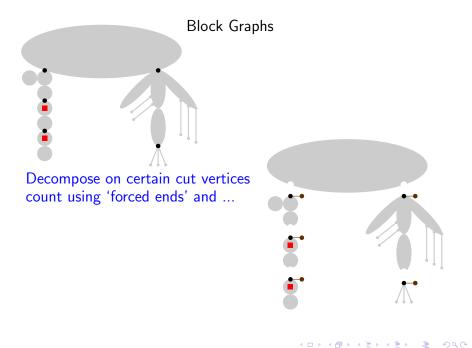
Block graphs, unit interval graphs with cut vertices, terminals T: complicated min = max ....

#### Linear Block Graphs (and unit interval with cut vertices)

Impose left right ordering and use 'forced ends' Decompose along carefully selected cut vertices i.e., those vertices do not pass through in left/right algorithm

#### Count the number of 'right' ends in each part

General block graphs: Similar idea with 'tops of paths' and parity of components 'below' special cut vertices



Trees with terminals T: Min |T|-Path Partition = Max C(G - U) - |U - T|

2-connected unit interval graphs with terminals T: Min |T|-Path Partition =  $\left\lceil \frac{|T|}{2} \right\rceil$ or +1 when ....

Threshold graphs with terminals T: Min |T|-Path Partition = Max  $\left(\sum \left\lceil \frac{|T_i|}{2} \right\rceil \right) + |R| - |U - T|$ or +1 when ....

Block graphs, unit interval graphs with cut vertices, terminals T: complicated min = max ....