## Interval Orders with Length Bounds

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## Intervals

## Interval Order



## Interval Graph

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### Interval Order??





Theorem (Wiener 1914) Interval Order  $\Leftrightarrow$  NO 2 + 2 R. Duncan Luce Example (1956) from Econometrica

- Add one grain of sugar at a time
- Cannot distinguish between consecutive cups

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- Can distinguish first and last
- Indifference is not transitive!
- Intervals with unit length (semiorders)



## Unit Interval Order??





Unit Interval Order  $\Leftrightarrow$  NO 2 + 2 & NO 3 + 1

Does a given order have a representation subject to given length bounds?

- Is there an efficient algorithm? Certifying algorithm?
- Forbidden suborder characterization?
- Graphs? Comparability invariant if all lower bounds equal

Assume interval order: no 2+2



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e.g., lengths between 1 and 4  $\Leftrightarrow$  no 1 + 6



e.g., Finite list of forbidden suborders but .... For lengths between 2 and 3: NO  $x \sim^2 \prec^4 y$  or  $x \sim \prec^2 \sim^2 \prec^3 y$  or  $x \sim \prec^3 \sim^2 \prec^2 y$ 

NO finite list for irrational



# Finite list of forbidden suborders for lengths 0 to nInfinite list of forbidden suborders for lengths 1 to n

Model Interval Order using Inequalities

 $\begin{array}{ll} l_1 \leq r_1 & r_1 < l_2 & \text{not } r_2 < l_3 \\ l_2 \leq r_2 & r_1 < l_3 & \text{not } r_3 < l_2 \\ l_3 \leq r_3 \end{array}$ 

Rewrite as:



We can represent an an order with intervals ⇔ Particular system of inequalities has a solution

Extends to:

- Constraints on interval length
- Minimize number of distinct endpoints
- Minimize 'support' length (if all lengths non-trivial)
- Partial information on ordering

• .....

Lemma (Farkas' Lemma 1906)

A system of inequalities has a solution  $\Leftrightarrow$  it is not inconsistent

 $Ax \le b$  has a solution or yA = 0,  $y \ge 0$ , yb < 0 has solution

Theorem (LP duality)  $\max \{cx | Ax \le b\} = \min \{yb | yA \ge c, y \ge 0\}$