Mathematics of Round Robin Tournaments

Garth Isaak Lehigh University

Math Club

Several times each year the math club meets for lectures by guest speakers or Bethel math professors. Twice this year famous mathematicians came to speak following banquets. The club provides additional resources for students interested in learning more about different aspects of the math field.

Photos by Tom Claassen and Eldon Esau



1st row: Ronald Quiring; 2nd row: Arnold Wedel, Alan Deckert; 3rd row: Garth Isaac, Amy Deckert; Rannie Goering, Steve Goering; 4th row: Dave King, John Thiesen, Al Dobbendick, Ron Headings

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Round Robin Tournaments

Seven players: Rempel, Wedel, Krehbiel, Helrich, Lehman, Brenneman, Thimm compete in a round robin tournament



Score sequence (5,4,3,3,3,2,1) records number of wins

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What questions might a mathematician (or ...) ask?



How do we rank players? What scores are possible? Structures in tournaments? Applications?

- We need tools from basic mathematics
- Linear algebra (Bethel Fall 1979)
- Combinatorics and Graph Theory (winter 1983)

Elementary Linear Algebra Final Exam December 12, 1979 1. Solve the system of equations 5.2x -y + Z=1 3 x +y - 3Z=2 9x -4Z=5 a) by writing as AK=B and multiplying both sides by A-1. b) by Cramer's rule, c) by Graussian elimination. 2. Solve the system dr = 2x - 4y - 22 $\frac{dy}{dt} = -x + 2y - 2$ 1= -2x +2y - 2 subject to the initial conditions × (0)=1, y (0)=2, Z(0)=-1, 3. Consider a simple model of an economy in which there are three "goods"; steel, food and labor. The production of each good consumes a part of what was preduced the year before, and the economist's guestion is whether land at what rate) the economy can expand, Suppose a new unit of steel regaines . 4 unit of existing steel and . 5 unit of labor, a unit of food requires . I unit of food and i 7 unit of labor, and producing (or maintaining) a unit of labor needs .8 unit of food and I unit each of steel and labor. Write amateix equation Vo = AV, relating the impats po, to the outputs A, , fr, and P, . Explain why

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5. Find the inverse of the matrix [-1 3 2]. to solve the systems x + y + z = 5 and x + y + z = 8 -x + 3y + 2z = 2 -x + 3y + 2z = -2 2x + y + z = 1 2x + y + z = 1 2x + y + z = 16. Suppose A, Band C are nxn matrices and ABC = @ Express A lin terms of B and C. © What can you say about C-1? © What can you express BT in terms of A and C? © Can you express BT in terms of A and C? 7. Prove that S is a subspace of a vector space Vit if S is nonempty and routsBES for all a, BES and real numbers vy S. (8) You may substitute this question, arising from last night's Club talls, for any question above for full evedit, or do it for Define the arithmetic mean, geometric mean and harmonic two numbers, represent them on a semicircle, and shar line segments you used have the desired long ths,

Fall 1979 Elementary Linear Algebra - Bethel College Professor Richard Rempel

- Mark Carpenter Ph.D. Mech Eng, Carnegie Mellon; NASA
- Allen Daubendiek Ph.D. EE, Florida(?); engineer
- Amy Deckert, Civil engineer
- William Ewy, engineer at HP Hong Kong
- Mark D. Friesen, software consultant
- John D. Harder, masters in statistics WSU/KSU; Ford
- Ron Headings, MBA Indiana; senior manager P&G research,
- Garth Isaak, Ph.D. Operations Research, Rutgers; Lehigh U
- J. David King, office 101st airborne
- Don Klippenstien, MD, Minnesota; Moffit Cancer center
- Richard Neufeld, nurse anesthesiologist
- Layne Reusser, MD, KSU; cardiovascular consultants, Wichita
- William Schmidt, Ph.D. Computer Science, Iowa; IBM
- John Thiesen, MS history WSU; Bethel

ACCK INTERTERM COURSE

Etle: COMBINATORICS AND GRAPH THEORY

- Amoription: Directed graphs, trees, circuits, paths, basic combinatorics, generating functions. Emphasis on applications and on use of computer problem solutions.
- Gredit: 3 credit hours towards mathematics or computer science.
- Hze and place: At McPherson College, Harnly Hall, Room 208, 1:30-4:30 daily January 3-27. First meeting 1:30, Jan. 3.

rincipal instructor: Richard Rempel, Bethel College(283-2500).

trequisites: Some calculus, some linear algebra (at least familiarity with matrices and their use in solving systems of linear equations), and a computer programming language.

ext: APPLIED COMBINATORICS, by Alan Tucker, Wiley, 1980. (On sale at the McPherson College Bookstore).

> Counting methods Elementary graph concepts Graph models Graph circuits and graph coloring Trees and searching

Generating functions Recurrence relations Inclusion-exclusion Polya's enumeration formula

sding: Based on homeworks and two exams.

ppics:

ecial Feature: Alan Tucker of SUNY at Stony Brook, the author of our text, will be a guest lecturer in the course for two days (dates and other details still not worked out).

can be down as Fr. Ja Ib can be done as - Ia Ie Je Je Ie draw graph with arrows indicating seeing another at the library (Note that I saw Fart A and B saw Fart but Ful A didn't see each other so wern it in library et save time two creaks a contradiction in times use ladification graph because seeing each other in the library is not transitive. FA IF IC JF To ---- cannot be placed without a contradiction since I d must overlap IF, IL, IA but not IB or IE so remain gust one of I con I E would not elleviate the inconsistancy , lade that D is the thirt

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ELA 8-29 ex. 1200 calpay pread moat 7 meat # 2.50 4 x as breador fruit as meat foren 754 in A 100 x + 50 4 + 300 Z = 1200 10x + 15y + 75z = 2,50 X + 4 # + + = 0 System of 3 signillaroous linen eq - in 3 unknowns + an 1st - 1st equation Zni - Unknow auxi+ a12 ×2 + a13 + - + a1NXN = b azixi+ azix2 + azz *3 m. +azNXN b2 ami +amitz+amitz --+anity 1 m -> * Inconsistent - NO solution * all bon = 07 honogeneores > never inconsistent (0, ... 0) never trivial sin 4 blue books assign - read 1.1; 1.2 solve A 8-31

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'Review' a little linear algebra this will be the hardest part, I promise



Do these have nonnegative solutions?

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Do these have nonnegative solutions?

$$\begin{array}{c} x + y + 2z = 3 \\ 5x + 8y + 13z = 21 \\ x - y + z = 0 \end{array} \begin{array}{c} x + y + 2z = 13 \\ 5x + 8y + 13z = 21 \\ x - 3y - 3z = 1 \end{array}$$

$$x = 0, \ y = z = 1 \\ yes \end{array} \qquad \begin{array}{c} \text{no} \\ \text{Why not?} \end{array}$$



Multiply equations by (-2), 1, 2 respectively

Multiply equations by (-2), 1, 2 respectively Add resulting equations

Multiply equations by (-2), 1, 2 respectively Add resulting equations

Result is

5x + 0y + 3z = -3

Every solution has at least one of x, y, z negative

Farkas' Lemma

Either a linear system has a nonnegative solution OR

There are multipliers showing inconsistency

$$\begin{array}{rrrr} -2 & x + y + 2z &= 13 \\ 1 & 5x + 8y + 13z &= 21 \end{array} \Rightarrow 3x + 6y + 9z = -5 \end{array}$$

also called Lagrange multipliers or Shadow prices This hints at Linear programming - significant use in applications Major LP advances reported in Wall Street Journal, NY Times etc including when I was in linear algebra at Bethel

- start with score sequences
- quick look at tournament structures
- quick look at rankings
- return to score sequences

10, 10, 9, 8, 8, 7, 5, 2, 2, 2, 2, 1

10, 10, 9, 8, 8, 7, 5, 2, 2, 2, 2, 1

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Simplest algorithm: test ALL possible tournaments

• Can we check if everyone in the room checks cases by hand for an hour?

10, 10, 9, 8, 8, 7, 5, 2, 2, 2, 2, 1

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- Can we check if everyone in the room checks cases by hand for an hour? NO
- Can we check if everyone in the room uses a computer for an hour?

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- Can we check using the world's fastest supercomputer for an hour?

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- Can we check if everyone in the room checks cases by hand for an hour? NO
- Can we check if everyone in the room uses a computer for an hour? NO
- Can we check using the world's fastest supercomputer for an hour? YES
- But we need the full hour

Is the following a possible score sequence for a tournament with 12 players? 10, 10, 9, 8, 8, 7, 5, 2, 2, 2, 1

- $\bullet\,$ There are $2^{66}\approx7\times10^{19}$ possible tournaments
- Fastest computer cluster runs at 10 Petaflops: 10×10^{15} operations per second
- We can check a size 12 problem in about an hour with this supercomputer.

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- What is the largest size we can check with 1,000 of these supercomputers?
- 12 (its not a typo) but with 10,000 supercomputers we can check size 13

22, 22, 20, 20, 20, 20, 19, 19, 18, 16, 16, 13, 13, 10, 8, 6, 6, 6, 5, 4, 4, 4, 3, 3, 3



22, 22, 20, 20, 20, 20, 19, 19, 18, 16, 16, 13, 13, 10, 8, 6, 6, 6, 5, 4, 4, 4, 3, 3, 3

Try checking ALL possible tournaments?

22, 22, 20, 20, 20, 20, 19, 19, 18, 16, 16, 13, 13, 10, 8, 6, 6, 6, 5, 4, 4, 4, 3, 3, 3

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UNIVERSE-ALL computer:

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All of the atoms in the known universe checking a billion tournaments per second

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Still not done checking all possibilities for this instance

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Try checking ALL possible tournaments?

UNIVERSE-ALL computer:

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Still not done checking all possibilities for this instance Use mathematical tools to make the check faster

Quick look at some structures in tournaments



beats all in 2 steps



path through all players application - Travelling Salesperson problem

These exercises involve no numbers or equations

Exercise

Prove that every tournament has a King

King - beats every other player or beats someone who beats them

Exercise

Prove that every Tournament has a Hamiltonian path

Hamiltonian path - sequence Rempel beats Wedel beats Kreibhel ... including each player exactly once

Exercise

Which Bethel Chemistry Professor did Garth attend High School with in Pennsylvania?

For those of you who prefer not to prove any theorems on a Saturday morning



Exercise

Every tournament either dominant player or at least three Kings Prove it What is the best way to rank players in a tournament? Depends on what 'best' means (need to define terms)

- Rank based on score sequence (number of wins)
- Rank to minimize the number of upsets
- Rank on 'iterated scores'

Iterated scores: K, H, L all have 3 wins K beats: 3+2+1=6H beats: 3+2+1=6L beats: 5+2+1=8

- First iterated scores: (14, 11, 6, 6, 8, 6, 4)
- Normalize (so they sum to 1) and repeat ...
- \Rightarrow Eigenvector
- Google's Page Rank similar



Ranking to minimize upsets





Can be different from ranking by wins Independently 'discovered'

- feedback arc sets (electrical engineering)
- linear ordering problem (comparing economic sectors)
- acyclic subgraph problem mathematics

Ranking to minimize upsets How bad can it get?



- Can ranking 100 players get 35 all wrong?
- YES
- We don't know for 99 players
- linear programming bound: $3n 1 \lfloor \log_2 n \rfloor$ for *n* players ranked wrong
- Similar bad things with ranking by scores
- Current research by graduate student Matt Prudente and former student Darren Narayan

Two problems

Is a list of numbers a score sequence?

Can we rank a tournament with at most k upsets

• One is 'easy'

nice necessary and sufficient conditions

- fast algorithm
- nice certificate when answer is 'no'
- One is 'hard'
 - 'Probably' no fast algorithms
 - 'Probably' no nice theorems exist
 - Not 'we are too dumb to find them' but in our model of mathematics and computation they can not exist



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\$1 Million dollar question (One of 7 Clay millenium problems) Is P = NP?

- *P* = polynomial (fast efficient algorithm)
- NP = nondeterministic Turing machine polynomial
- NP can verify a solution from 'god' efficiently
- NP-complete: at least as 'hard' as other NP problems
- To show *P* = *NP* find a fast algorithm *NP*-complete problem e.g. for Travelling Salesperson Problem
- To show P ≠ NP prove that no efficient algorithm exists for NP-complete problems

• Ranking to minimize upsets is NP-complete

- 'Probably' no nice theorems can exist
- If $P \neq NP$ and problem is NP-complete no such theorem can exist
- Problems like these arise in applications. Use linear programming to get bounds and approximations
- Recognizing score sequences has nice theorems and fast algorithms

22, 22, 20, 20, 20, 20, 19, 19, 18, 16, 16, 13, 13, 10, 8, 6, 6, 6, 5, 4, 4, 4, 3, 3, 3

- Want necessary and sufficient conditions
 - We heard about necessary and sufficient conditions in yesterday's talk
- Can certify 'YES' by producing a tournament with these scores
- Want 'nice' certificate is answer is NO



Which of the following are score sequences for a tournament with 7 players?

 $(7, 5, 4\frac{1}{3}, 4, 2\frac{3}{7}, 0, -2)$ (5, 4, 3, 3, 3, 1, 0) (3, 3, 3, 3, 3, 3, 3, 3) (6, 6, 4, 2, 1, 1, 1)

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Which of the following are score sequences for a tournament with 7 players?

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(6, 6, 4, 2, 1, 1, 1) NO - two teams cannot win all of their games

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- Can the two lowest scores both be 0?
- Can the three lowest scores be 1, 1, 0?
- Can the four lowest scores be 2, 1, 1, 1?

- Can the two lowest scores both be 0? NO - they play each other
- Can the three lowest scores be 1,1,0? NO - they play 3 games
- Can the four lowest scores be 2, 1, 1, 1? NO - they play 6 games

• :

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- Seven teams play $21 = \frac{7(7-1)}{2}$ games so the lowest seven scores must sum to at least 21
- k teams play $\frac{k(k-1)}{2}$ games so they have at least that many wins

Landau (1951) considered tournaments in the context of pecking order in poultry populations

Necessary condition

The number of wins for any set of teams must be as large as the number of games played between those teams

Landau's Theorem:

This necessary condition is also sufficient

That is: Not a score sequence $\,\Rightarrow\,$ there is a set of teams violating these obvious conditions

The sequence 22, 22, 20, 20, 20, 20, 19, 19, 18, 16, 16, 13, 13, 10, 8, 6, 6, 6, 5, 4, 4, 4, 3, 3, 3 can be checked by hand in a few minutes. It is not a score sequence

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The sequence 22, 22, 20, 20, 20, 20, 19, 19, 18, 16, 16, 13, 13, 10, 8, 6, 6, 6, 5, 4, 4, 4, 3, 3, 3 can be checked by hand in a few minutes. It is not a score sequence

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Not a score sequence

Last 10 teams have 44 wins in $45 = \frac{10.9}{2}$ games

$$\sum_{i\in I} s_i \geq \binom{|I|}{2}$$
 for any $I \subseteq \{1, 2, \dots, n\}$

with equality when $I = \{1, 2, \ldots, n\}$

There are many proofs: by induction ...

Try modeling score sequences using equations Shadow prices (Lagrange multipliers) plus ... yield a proof

$$\sum_{i \in I} s_i \ge \binom{|I|}{2} \text{ for any } I \subseteq \{1, 2, \dots, n\}$$

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What if we allow ties?

What if the score is 3 points for a win, 1 for a tie and 0 for a loss (world cup soccer scoring)?

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What if we allow ties?

This problem is not solved Summer 2010 REU, Amy, Hannah, Jenny used a discrete tomography problem to show its NP-hard for boy/girl version

What if the score is 3 points for a win, 1 for a tie and 0 for a loss (world cup soccer scoring)? This problem is not solved







 Ron Headings helps Herman Bubbert with his homework.

2. Arnold Wedel, Math Professor. 3. Richard Rempel, Math Professor.

Thanks to Bethel Faculty!