

This is closed book, closed notes etc.

You have 50 minutes to take this exam.

Points for each problem are indicated as [·].

Be sure to do the problems that you know well first.

There are 6 problems plus a bonus problem

When you are asked to pick a part or parts of a problem to answer clearly indicate your selection. If you attempt more than is asked those with lowest scores will be used.

**1:** [25] Pick *two* of the following and for each give a proof by *induction*.

(a)  $\sum_{i=1}^n i = \frac{n(n+1)}{2}$  for  $n = 1, 2, \dots$

(b) The Fibonacci numbers  $F_n$  given by  $F_n = F_{n-1} + F_{n-2}$  for  $n \geq 2$  with  $F_0 = 0$  and  $F_1 = 1$  satisfy  $1 + \sum_{i=0}^n F_i = F_{n+2}$

(c) If  $D_n = D_{n-1} + 12D_{n-2}$  for  $n \geq 2$  with  $D_0 = 3$  and  $D_1 = -2$  then  $D_n = 2 \cdot (-3)^n + 4^n$ .

**2:** [17] Pick *two* of the following and give combinatorial (in terms of sets) proofs. You should just use the interpretation of  $\binom{n}{k}$  as the number of size  $k$  subsets of an  $n$  set. You should not write any formula using  $\binom{n}{k} = \frac{n!}{(n-k)!k!}$ .

(a) (Pascal's identity)  $\binom{n}{k} = \binom{n-1}{k-1} + \binom{n-1}{k}$

(b)  $\binom{n}{k} = \binom{n}{n-k}$

(c)  $k \binom{n}{k} = n \binom{n-1}{k-1}$

**3:** [11] Prove the Binomial Theorem  $(x + y)^n = \sum_{k=0}^n \binom{n}{k} x^k y^{n-k}$  in any way you like.

**4:** [13] Make use of the binomial theorem to prove that the number of subsets of an  $n$  element set is  $2^n$  and that the number of even sized subsets of an  $n$  element set is equal to the number of odd sized subsets.

**5:** [15] Make use of  $n^3 = \sum_{i=1}^n (i^3 - (i-1)^3)$  to show that  $\sum_{i=1}^n i^2 = \frac{n(n+1)(2n+1)}{6}$ . You may use  $\sum_{i=1}^n i = \frac{n(n+1)}{2}$ .

**6:** [19] Consider a bipartite graph  $G$  with parts  $X, Y$  such that  $|X| = |Y|$ .

Hall's condition is  $|T| \leq |N(T)|$  for all  $T \subseteq X$

The marriage condition is  $|C| \geq n$  for every vertex cover. Each is a necessary and sufficient condition for  $G$  to have a perfect matching. Consider the following four statements:

- (1) Hall's condition is a necessary condition for  $G$  to have a perfect matching
- (2) Hall's condition is a sufficient condition for  $G$  to have a perfect matching
- (3) The marriage condition is a necessary condition for  $G$  to have a perfect matching
- (4) The marriage condition is a sufficient condition for  $G$  to have a perfect matching

(a) Which one of (1) or (2) is the same as: If  $G$  does not have a perfect matching then there is a  $T \subseteq X$  such that  $|T| > |N(T)|$ ?

(b) Which one of (1) or (2) is the same as: if  $G$  has a perfect matching then  $|T| \leq |N(T)|$  for all  $T \subseteq X$ ?

(c) State the contrapositive of the statement in (a)

(d) State the contrapositive of the statement in (b)

(e) State a condition and its contrapositive for (3) analogous to those described in (a) and (b)

(f) State a condition and its contrapositive for (4) analogous to those described in (a) and (b)

**7:** [16] Pick *one* of the following and prove it:

**(a)** (If  $G$  does not have a perfect matching then there is a  $T \subseteq X$  such that  $|T| > |N(T)|$ ) implies (If  $G$  does not have a perfect matching then  $G$  has a vertex cover  $C$  with  $|C| < n$ ).

**(b)** (If  $G$  does not have a perfect matching then  $G$  has a vertex cover  $C$  with  $|C| < n$ ) implies (If  $G$  does not have a perfect matching then there is a  $T \subseteq X$  such that  $|T| > |N(T)|$ )